# Compiler Construction

ICOM 4029

1:30-3:10

CID 201

#### ICOM 4029 - Outline

- Prontuario
- · Course Outline
- Brief History of PLs
- · Programming Language Design Criteria
- · Programming Language Implementation

## Programming Assignments Highlights

- Implement a compiler in four phases
- Teams of two students (Choose your partner!)
- Development in C++ or Java for Linux
- Use Academic Computer Center (Amadeus) if needed
- Can work on your personal computers
- Source Language = COOL (UC Berkeley CS164)
- Target Language = MIPS Assembly (SPIM)
- Each project must have some unique feature chosen by the development team
- Each compiler must pass a minimal set of tests in order to pass the class.

#### Homework for next week

- Read the COOL Reference Manual
- Choose your partner
  - notify me by email
- Choose your development language
  - C++ or Java
- Read the Flex (C++) or JLex (Java) manual

# (Short) History of High-Level Languages

- 1953 IBM develops the 701
- All programming done in assembly
- Problem: Software costs exceeded hardware costs!
- · John Backus: "Speedcoding"
  - An interpreter
  - Ran 10-20 times slower than hand-written assembly

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#### FORTRAN I

- 1954 IBM develops the 704
- John Backus
  - Idea: translate high-level code to assembly
  - Many thought this impossible
    - · Had already failed in other projects
- 1954-7 FORTRAN I project
- By 1958, >50% of all software is in FORTRAN
- Cut development time dramatically
  - $(2 \text{ wks} \rightarrow 2 \text{ hrs})$

#### FORTRAN I

- The first compiler
  - Produced code almost as good as hand-written
  - Huge impact on computer science
- · Led to an enormous body of theoretical work
- Modern compilers preserve the outlines of FORTRAN I

#### History of Ideas: Abstraction

- Abstraction = detached from concrete details
- Abstraction necessary to build software systems
- Modes of abstraction
  - Via languages/compilers:
    - · Higher-level code, few machine dependencies
  - Via subroutines
    - Abstract interface to behavior
  - Via modules
    - Export interfaces; hide implementation
  - Via abstract data types
    - Bundle data with its operations

### History of Ideas: Types

- Originally, few types
  - FORTRAN: scalars, arrays
  - LISP: no static type distinctions
- · Realization: Types help
  - Allow the programmer to express abstraction
  - Allow the compiler to check against many frequent errors
  - Sometimes to the point that programs are guaranteed "safe"
- More recently
  - Lots of interest in types
  - Experiments with various forms of parameterization
  - Best developed in functional programming

#### History of Ideas: Reuse

- Reuse = exploits common patterns in software systems
- Goal: mass-produced software components
- Reuse is difficult
- Two popular approaches (combined in C++)
  - Type parameterization (List(int), List(double))
  - Classes and inheritance: C++ derived classes
- Inheritance allows
  - Specialization of existing abstraction
  - Extension, modification, hiding behavior

## Programming Language Economics 101

- · Languages are adopted to fill a void
  - Enable a previously difficult/impossible application
  - Orthogonal to language design quality (almost)
- · Programmer training is the dominant cost
  - Languages with many users are replaced rarely
  - Popular languages become ossified
  - But easy to start in a new niche . . .

# Why So Many Languages?

- Application domains have distinctive (and conflicting) needs
- Examples:
  - Scientific Computing: high performance
  - Business: report generation
  - Artificial intelligence: symbolic computation
  - Systems programming: low-level access
  - Special purpose languages

### Topic: Language Design

- No universally accepted metrics for design
- "A good language is one people use"?
- · NO!
  - Is COBOL the best language?
- · Good language design is hard

# Language Evaluation Criteria

Characteristic	Criteria		
	Readability	Writeability	Reliability
Simplicity	*	*	*
Data types	*	*	*
Syntax design	*	*	*
Abstraction		*	*
Expressivity		*	*
Type checking			*
Exception handling			*

# Why Study Languages and Compilers?

- Increase capacity of expression
- Improve understanding of program behavior
- Increase ability to learn new languages
- Learn to build a large and reliable system
- · See many basic CS concepts at work

#### **Trends**

- Language design
  - Many new special-purpose languages
  - Popular languages to stay
- Compilers
  - More needed and more complex
  - Driven by increasing gap between
    - new languages
    - new architectures
  - Venerable and healthy area

# How are Languages Implemented?

- Two major strategies:
  - Interpreters (older, less studied)
  - Compilers (newer, much more studied)
- · Interpreters run programs "as is"
  - Little or no preprocessing
- Compilers do extensive preprocessing

# Language Implementations

- Batch compilation systems dominate
  - E.g., gcc
- Some languages are primarily interpreted
  - E.g., Java bytecode
- Some environments (Lisp) provide both
  - Interpreter for development
  - Compiler for production

# The Structure of a Compiler

- 1. Lexical Analysis
- 2. Parsing
- 3. Semantic Analysis
- 4. Optimization
- 5. Code Generation

The first 3, at least, can be understood by analogy to how humans comprehend English.

## Lexical Analysis

- First step: recognize words.
  - Smallest unit above letters

#### This is a sentence.

- · Note the
  - Capital "T" (start of sentence symbol)
  - Blank " " (word separator)
  - Period "." (end of sentence symbol)

# More Lexical Analysis

Lexical analysis is not trivial. Consider:

ist his ase nte nce

 Plus, programming languages are typically more cryptic than English:

$$p\rightarrow f ++ = -.12345e-5$$

#### And More Lexical Analysis

 Lexical analyzer divides program text into "words" or "tokens"

if 
$$x == y$$
 then  $z = 1$ ; else  $z = 2$ ;

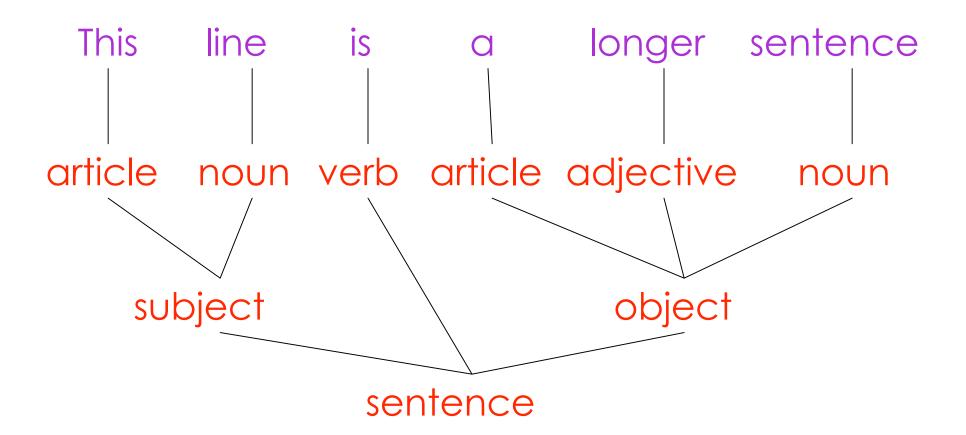
· Units:

if, 
$$x$$
, ==,  $y$ , then,  $z$ , =, 1, ;, else,  $z$ , =, 2, ;

### Parsing

- Once words are understood, the next step is to understand sentence structure
- Parsing = Diagramming Sentences
  - The diagram is a tree

# Diagramming a Sentence

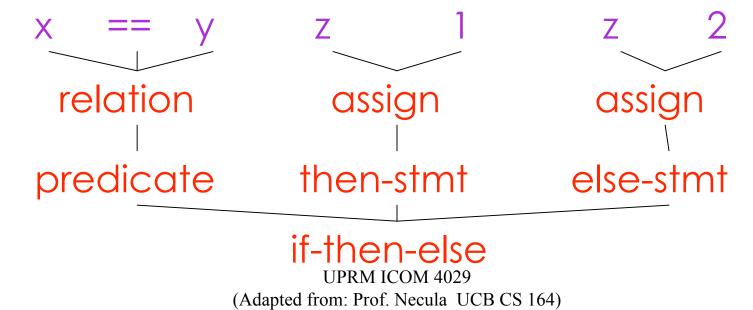


## Parsing Programs

- Parsing program expressions is the same
- · Consider:

If 
$$x == y$$
 then  $z = 1$ ; else  $z = 2$ ;

Diagrammed:



# Semantic Analysis

- Once sentence structure is understood, we can try to understand "meaning"
  - But meaning is too hard for compilers
- Compilers perform limited analysis to catch inconsistencies
- Some do more analysis to improve the performance of the program

## Semantic Analysis in English

#### Example:

Jack said Jerry left his assignment at home. What does "his" refer to? Jack or Jerry?

#### · Even worse:

Jack said Jack left his assignment at home?

How many Jacks are there?

Which one left the assignment?

# Semantic Analysis in Programming

- Programming languages define strict rules to avoid such ambiguities
- This C++ code prints "4"; the inner definition is used

```
{
  int Jack = 3;
  {
    int Jack = 4;
    cout << Jack;
}</pre>
```

## More Semantic Analysis

 Compilers perform many semantic checks besides variable bindings

• Example:

Jack left her homework at home.

- A "type mismatch" between her and Jack; we know they are different people
  - Presumably Jack is male

#### Examples of Semantic Checks in PLs

- Variables defined before used
- Variables defined once
- Type compatibility
- Correct arguments to functions
- · Constants are not modified
- · Inheritance hierarchy has no cycles

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### Optimization

- No strong counterpart in English, but akin to editing
- Automatically modify programs so that they
  - Run faster
  - Use less memory
  - In general, conserve some resource
- The project has no optimization component

# Optimization Example

X = Y \* 0 is the same as X = 0

#### NO!

Valid for integers, but not for floating point numbers

## Examples of common optimizations in PLs

- Dead code elimination
- Evaluating repeated expressions only once
- Replace expressions by simpler equivalent expressions
- Evaluate expressions at compile time
- Inline procedures
- Move constant expressions out of loops

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#### Code Generation

- Produces assembly code (usually)
- A translation into another language
  - Analogous to human translation

## Intermediate Languages

- Many compilers perform translations between successive intermediate forms
  - All but first and last are *intermediate languages* internal to the compiler
  - Typically there is 1 IL
- IL's generally ordered in descending level of abstraction
  - Highest is source
  - Lowest is assembly

# Intermediate Languages (Cont.)

- IL's are useful because lower levels expose features hidden by higher levels
  - registers
  - memory layout
  - etc.
- But lower levels obscure high-level meaning

#### **Issues**

- Compiling is almost this simple, but there are many pitfalls.
- Example: How are erroneous programs handled?
- Language design has big impact on compiler
  - Determines what is easy and hard to compile
  - Course theme: many trade-offs in language design

# Compilers Today

- The overall structure of almost every compiler adheres to our outline
- The proportions have changed since FORTRAN
  - Early: lexing, parsing most complex, expensive
  - Today: optimization dominates all other phases, lexing and parsing are cheap

## Trends in Compilation

- · Compilation for speed is less interesting. But:
  - scientific programs
  - advanced processors (Digital Signal Processors, advanced speculative architectures)
- Ideas from compilation used for improving code reliability:
  - memory safety
  - detecting concurrency errors (data races)
  - ...