Compiler Construction

ICOM 4029

7:30-9:15 PM

5-227

ICOM 4029 - Outline

- Prontuario
- · Course Outline
- Brief History of PLs
- Programming Language Design Criteria
- · Programming Language Implementation

Programming Assignments Highlights

- Implement a compiler in four phases
- Teams of two students (Choose your partner!)
- Development in Java
- Use Academic Computer Center (Amadeus) if needed
- Can work on your personal computers
- Source Language = COOL (UC Berkeley CS164)
- Target Language = MIPS Assembly (SPIM)
- Each compiler must pass a minimal set of tests in order to pass the class.

Homework for next week

- Read the COOL Reference Manual
- Choose your partner
 - notify me by email
- Read the JLex (Java) manual

(Short) History of High-Level Languages

- 1953 IBM develops the 701
- · All programming done in assembly
- Problem: Software costs exceeded hardware costs!
- · John Backus: "Speedcoding"
 - An interpreter
 - Ran 10-20 times slower than hand-written assembly

FORTRAN I

- 1954 IBM develops the 704
- John Backus
 - Idea: translate high-level code to assembly
 - Many thought this impossible
 - Had already failed in other projects
- 1954-7 FORTRAN I project
- By 1958, >50% of all software is in FORTRAN
- Cut development time dramatically
 - $(2 \text{ wks} \rightarrow 2 \text{ hrs})$

FORTRAN I

- The first compiler
 - Produced code almost as good as hand-written
 - Huge impact on computer science
- · Led to an enormous body of theoretical work
- Modern compilers preserve the outlines of FORTRAN I

History of Ideas: Abstraction

- · Abstraction = detached from concrete details
- Abstraction necessary to build software systems
- Modes of abstraction
 - Via languages/compilers:
 - · Higher-level code, few machine dependencies
 - Via subroutines
 - Abstract interface to behavior
 - Via modules
 - · Export interfaces; hide implementation
 - Via abstract data types
 - Bundle data with its operations

History of Ideas: Types

- Originally, few types
 - FORTRAN: scalars, arrays
 - LISP: no static type distinctions
- · Realization: Types help
 - Allow the programmer to express abstraction
 - Allow the compiler to check against many frequent errors
 - Sometimes to the point that programs are guaranteed "safe"
- · More recently
 - Lots of interest in types
 - Experiments with various forms of parameterization
 - Best developed in functional programming

History of Ideas: Reuse

- Reuse = exploits common patterns in software systems
- · Goal: mass-produced software components
- Reuse is difficult
- Two popular approaches (combined in C++)
 - Type parameterization (List(int), List(double))
 - Classes and inheritance: C++ derived classes
- Inheritance allows
 - Specialization of existing abstraction
 - Extension, modification, hiding behavior

Programming Language Economics 101

- · Languages are adopted to fill a void
 - Enable a previously difficult/impossible application
 - Orthogonal to language design quality (almost)
- Programmer training is the dominant cost
 - Languages with many users are replaced rarely
 - Popular languages become ossified
 - But easy to start in a new niche . . .

Why So Many Languages?

- Application domains have distinctive (and conflicting) needs
- Examples:
 - Scientific Computing: high performance
 - Business: report generation
 - Artificial intelligence: symbolic computation
 - Systems programming: low-level access
 - Special purpose languages

Topic: Language Design

- No universally accepted metrics for design
- "A good language is one people use"?
- · NO!
 - Is COBOL the best language?
- Good language design is hard

Language Evaluation Criteria

Characteristic	Criteria		
	Readability	Writeability	Reliability
Simplicity	*	*	*
Data types	*	*	*
Syntax design	*	*	*
Abstraction		*	*
Expressivity		*	*
Type checking			*
Exception handling			*

Why Study Languages and Compilers?

- Increase capacity of expression
- Improve understanding of program behavior
- Increase ability to learn new languages
- · Learn to build a large and reliable system
- See many basic CS concepts at work

Trends

- Language design
 - Many new special-purpose languages
 - Popular languages to stay
- Compilers
 - More needed and more complex
 - Driven by increasing gap between
 - new languages
 - · new architectures
 - Venerable and healthy area

How are Languages Implemented?

- Two major strategies:
 - Interpreters (older, less studied)
 - Compilers (newer, much more studied)
- Interpreters run programs "as is"
 - Little or no preprocessing
- · Compilers do extensive preprocessing

Language Implementations

- Batch compilation systems dominate
 - E.g., gcc
- Some languages are primarily interpreted
 - E.g., Java bytecode
- · Some environments (Lisp) provide both
 - Interpreter for development
 - Compiler for production

The Structure of a Compiler

- 1. Lexical Analysis
- 2. Parsing
- 3. Semantic Analysis
- 4. Optimization
- 5. Code Generation

The first 3, at least, can be understood by analogy to how humans comprehend English.

Lexical Analysis

- First step: recognize words.
 - Smallest unit above letters

This is a sentence.

- · Note the
 - Capital "T" (start of sentence symbol)
 - Blank " " (word separator)
 - Period "." (end of sentence symbol)

More Lexical Analysis

· Lexical analysis is not trivial. Consider:

ist his ase nte nce

 Plus, programming languages are typically more cryptic than English:

$$p\rightarrow f ++ = -.12345e-5$$

And More Lexical Analysis

 Lexical analyzer divides program text into "words" or "tokens"

if
$$x == y$$
 then $z = 1$; else $z = 2$;

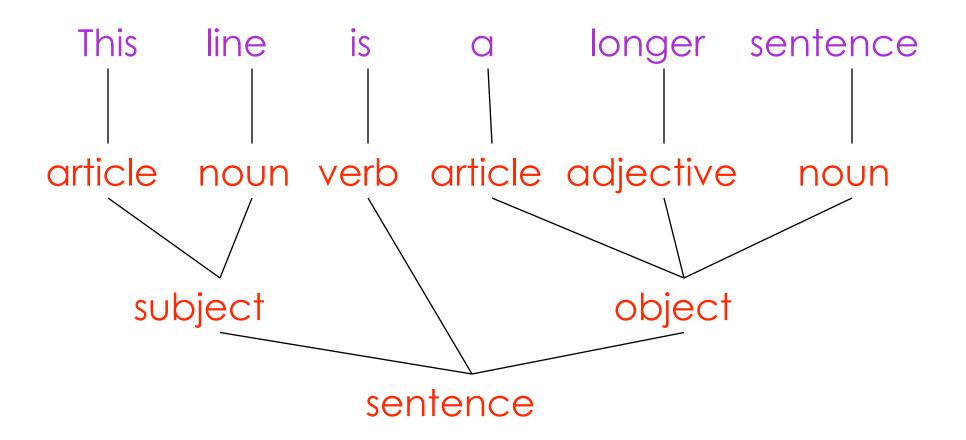
· Units:

if,
$$x$$
, ==, y , then, z , =, 1, ;, else, z , =, 2, ;

Parsing

- Once words are understood, the next step is to understand sentence structure
- Parsing = Diagramming Sentences
 - The diagram is a tree

Diagramming a Sentence

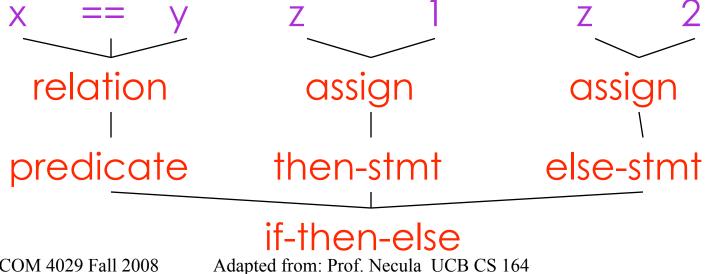


Parsing Programs

- Parsing program expressions is the same
- · Consider:

If
$$x == y$$
 then $z = 1$; else $z = 2$;

Diagrammed:



Semantic Analysis

- Once sentence structure is understood, we can try to understand "meaning"
 - But meaning is too hard for compilers
- Compilers perform limited analysis to catch inconsistencies
- Some do more analysis to improve the performance of the program

Semantic Analysis in English

· Example:

Jack said Jerry left his assignment at home. What does "his" refer to? Jack or Jerry?

• Even worse:

Jack said Jack left his assignment at home?

How many Jacks are there?

Which one left the assignment?

Semantic Analysis in Programming

- Programming languages define strict rules to avoid such ambiguities
- This C++ code prints "4"; the inner definition is used

```
{
  int Jack = 3;
  {
    int Jack = 4;
    cout << Jack;
}
</pre>
```

More Semantic Analysis

 Compilers perform many semantic checks besides variable bindings

Example:

Jack left her homework at home.

- A "type mismatch" between her and Jack; we know they are different people
 - Presumably Jack is male

Examples of Semantic Checks in PLs

- Variables defined before used
- Variables defined once
- Type compatibility
- Correct arguments to functions
- · Constants are not modified
- Inheritance hierarchy has no cycles

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Optimization

- No strong counterpart in English, but akin to editing
- Automatically modify programs so that they
 - Run faster
 - Use less memory
 - In general, conserve some resource
- The project has no optimization component

Optimization Example

X = Y * 0 is the same as X = 0

NO!

Valid for integers, but not for floating point numbers

Examples of common optimizations in PLs

- Dead code elimination
- · Evaluating repeated expressions only once
- Replace expressions by simpler equivalent expressions
- · Evaluate expressions at compile time
- Inline procedures
- Move constant expressions out of loops

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Code Generation

- Produces assembly code (usually)
- A translation into another language
 - Analogous to human translation

Intermediate Languages

- Many compilers perform translations between successive intermediate forms
 - All but first and last are intermediate languages internal to the compiler
 - Typically there is 1 IL
- IL's generally ordered in descending level of abstraction
 - Highest is source
 - Lowest is assembly

Intermediate Languages (Cont.)

- IL's are useful because lower levels expose features hidden by higher levels
 - registers
 - memory layout
 - etc.
- · But lower levels obscure high-level meaning

Issues

- Compiling is almost this simple, but there are many pitfalls.
- Example: How are erroneous programs handled?
- · Language design has big impact on compiler
 - Determines what is easy and hard to compile
 - Course theme: many trade-offs in language design

Compilers Today

- The overall structure of almost every compiler adheres to our outline
- The proportions have changed since FORTRAN
 - Early: lexing, parsing most complex, expensive
 - Today: optimization dominates all other phases, lexing and parsing are cheap

Trends in Compilation

- · Compilation for speed is less interesting. But:
 - scientific programs
 - advanced processors (Digital Signal Processors, advanced speculative architectures)
- Ideas from compilation used for improving code reliability:
 - memory safety
 - detecting concurrency errors (data races)
 - ...