

ICOM 4075:
Foundations of Computing

Lecture 6
Functions (2)

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Lecture Notes Originally Written By Prof. Yi Qian

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- Section 2.1: (pp.88-90)

1. b. d.
2. b. d. f.
3. b. d.
4. b.
6. b. d.
7. b. d. f.
8. b.
- 9.
- 10.
11. b.
13. b. d.
14. b. d.
15. b. d. f.
- 17.

Reading

- Textbook: James L. Hein, *Discrete Structures, Logic, and Computability*, 2nd edition, Chapter 2. Section 2.2

Constructing Functions

- Composition of Functions
 - *Composition* of functions is a natural process that we often use without even thinking.
 - E.g., $\text{floor}(\log_2(6))$ involves the composition of the two functions floor and \log_2 . To evaluate the expression, we first evaluate $\log_2(6)$, which is a number between 2 and 3. Then we apply the floor function to this number, obtaining the value 2.
- Definition of Composition
 - The *composition* of two functions f and g is the function denoted by $f \circ g$ and defined by $(f \circ g)(x) = f(g(x))$.
- Notice that composition makes sense only for values of x in the domain of g such that $g(x)$ is in the domain of f .
 - So if $g: A \rightarrow B$ and $f: C \rightarrow D$ and $B \subset C$, then the composition $f \circ g$ makes sense. In other words, for every $x \in A$ it follows that $g(x) \in B$, and since $B \subset C$ it follows that $f(g(x)) \in D$. It also follows that $f \circ g: A \rightarrow D$.
 - E.g., $\log_2: \mathbb{R}^+ \rightarrow \mathbb{R}$ and $\text{floor}: \mathbb{R} \rightarrow \mathbb{Z}$, where \mathbb{R}^+ denotes the set of positive real numbers. So for any positive real number x , the expression $\log_2(x)$ is a real number and thus $\text{floor}(\log_2(x))$ is an integer. So the composition $\text{floor} \circ \log_2: \mathbb{R}^+ \rightarrow \mathbb{Z}$.

Composition of Functions

- Composition of functions is *associative*:
 - If f , g , and h are functions of the right type such that $(f \circ g) \circ h$ and $f \circ (g \circ h)$ make sense, then $(f \circ g) \circ h = f \circ (g \circ h)$.
- Composition of functions is *not commutative*:
 - E.g., suppose that f and g are defined by $f(x) = x + 1$ and $g(x) = x^2$. To show that $f \circ g \neq g \circ f$, we only need to find one number x such that $(f \circ g)(x) \neq (g \circ f)(x)$. We'll try $x = 3$ and observe that
$$(f \circ g)(3) = f(g(3)) = f(3^2) = 3^2 + 1 = 10.$$
$$(g \circ f)(3) = g(f(3)) = g(3 + 1) = (3 + 1)^2 = 16.$$
Therefore, $(f \circ g)(3) \neq (g \circ f)(3)$
- A function that always returns its argument is called an *identity* function. For a set A we sometimes write “ id_A ” to denote the identity function defined by $\text{id}_A(a) = a$ for all $a \in A$. If $f: A \rightarrow B$, then we always have the following equation: $f \circ \text{id}_A = f = \text{id}_B \circ f$

The Sequence, Distribute, and Pairs Functions

- The *sequence* function “seq” has type $N \rightarrow \text{lists}(N)$ and is defined as follows for any natural number n : $\text{seq}(n) = \langle 0, 1, \dots, n \rangle$.
 - E.g., $\text{seq}(0) = \langle 0 \rangle$, $\text{seq}(2) = \langle 0, 1, 2 \rangle$, $\text{seq}(5) = \langle 0, 1, 2, 3, 4, 5 \rangle$.
- The *distribute* function “dist” has type $A \times \text{lists}(B) \rightarrow \text{lists}(A \times B)$. It takes an element x from A and a list y from $\text{lists}(B)$ and returns the list of pairs made up by pairing x with each element of y .
 - E.g., $\text{dist}(x, \langle r, s, t \rangle) = \langle (x, r), (x, s), (x, t) \rangle$.
- The *pairs* function takes two lists of equal length and returns the list of pairs of corresponding elements.
 - E.g., $\text{pairs}(\langle a, b, c \rangle, \langle d, e, f \rangle) = \langle (a, d), (b, e), (c, f) \rangle$.
 - Since the domain of pairs is a proper subset of $\text{lists}(A) \times \text{lists}(B)$, it is a partial function of type $\text{lists}(A) \times \text{lists}(B) \rightarrow \text{lists}(A \times B)$.

Composing Functions with Different Arities

- Composition can also occur between functions with different arities.
 - E.g., suppose we define the following function $f(x, y) = \text{dist}(x, \text{seq}(y))$. In this case dist has two arguments and seq has one argument. For example, we'll evaluate the expression $f(5, 3)$.

$$\begin{aligned} f(5, 3) &= \text{dist}(5, \text{seq}(3)) \\ &= \text{dist}(5, \langle 0, 1, 2, 3 \rangle) \\ &= \langle (5, 0), (5, 1), (5, 2), (5, 3) \rangle. \end{aligned}$$

Distribute a Sequence

- We'll show that the definition $f(x, y) = \text{dist}(x, \text{seq}(y))$ is a special case of the following more general form of *composition*, where X can be replaced by any number of arguments. $f(X) = h(g_1(X), \dots, g_n(X))$.

- Distribute a Sequence

- We'll show that the definition $f(x, y) = \text{dist}(x, \text{seq}(y))$ fits the general form of composition. To make it fit the form, we'll define the functions $\text{one}(x, y) = x$ and $\text{two}(x, y) = y$. Then we have the following representation of f .

$$\begin{aligned} f(x, y) &= \text{dist}(x, \text{seq}(y)) \\ &= \text{dist}(\text{one}(x, y), \text{seq}(\text{two}(x, y))) \\ &= \text{dist}(\text{one}(x, y), (\text{seq} \circ \text{two}(x, y))). \end{aligned}$$

The last expression has the general form of composition

$$f(X) = h(g_1(X), g_2(X)),$$

where $X = (x, y)$, $h = \text{dist}$, $g_1 = \text{one}$, and $g_2 = \text{seq} \circ \text{two}$

The Max Function

- The Max Function

- Suppose we define the function “max”, to return the maximum of two numbers as follows:

$$\text{max}(x, y) = \text{if } x < y \text{ then } y \text{ else } x.$$







Then we can use max to define the function “max3”, which returns the maximum of three numbers, by the following composition:

$$\text{max3}(x, y, z) = \text{max}(\text{max}(x, y), z).$$

Minimum Depth of a Binary Tree

- To find the minimum depth of a binary tree in terms of the numbers of nodes:
 - The following figure lists a few sample cases in which the trees are as compact as possible, which means that they have the least depth for the number of nodes. Let n denote the number of nodes. Notice that when $4 \leq n < 8$, the depth is 2. Similarly, the depth is 3 whenever $8 \leq n < 16$.
 - At the same time we know that $\log_2(4) = 2$, $\log_2(8) = 3$, and for $4 \leq n < 8$ we have $2 \leq \log_2(n) < 3$. So $\log_2(n)$ almost works as the depth function.
 - In general, we have the minimum depth function as the composition of the floor function and the \log_2 function:

$$\text{minDepth}(n) = \text{floor}(\log_2(n)).$$

Binary tree	Nodes	Depth
	1	0
	2	1
	3	1
	4	2
	7	2
	15	3

List of Pairs

- Suppose we want to construct a definition for the following function in terms of known functions

$$f(n) = \langle (0, 0), (1, 1), \dots, (n, n) \rangle \text{ for any } n \in \mathbb{N}.$$

Starting with this informal definition, we'll transform it into a composition of known functions.

$$\begin{aligned} f(n) &= \langle (0, 0), (1, 1), \dots, (n, n) \rangle \\ &= \text{pairs}(\langle 0, 1, \dots, n \rangle, \langle 0, 1, \dots, n \rangle) \\ &= \text{pairs}(\text{seq}(n), \text{seq}(n)). \end{aligned}$$

- Suppose we want to construct a definition for the following function in terms of known functions

$$g(k) = \langle (k, 0), (k, 1), \dots, (k, k) \rangle \text{ for any } k \in \mathbb{N}.$$

Starting with this informal definition, we'll transform it into a composition of known functions.

$$\begin{aligned} g(k) &= \langle (k, 0), (k, 1), \dots, (k, k) \rangle \\ &= \text{dist}(k, \langle 0, 1, \dots, k \rangle) \\ &= \text{dist}(k, \text{seq}(k)). \end{aligned}$$

The Map Function

- Definition of the Map Function:

- Let f be a function with domain A and let $\langle x_1, \dots, x_n \rangle$ be a list of elements from A . Then

$$\text{map}(f, \langle x_1, \dots, x_n \rangle) = \langle f(x_1), \dots, f(x_n) \rangle.$$

So the type of the map function can be written as

$$\text{map}: (A \rightarrow B) \times \text{list}(A) \rightarrow \text{lists}(B).$$

- E.g.,

$$\text{map}(\text{floor}, \langle -1.5, -0.5, 0.5, 1.5, 2.5 \rangle)$$

$$= \langle \text{floor}(-1.5), \text{floor}(-0.5), \text{floor}(0.5), \text{floor}(1.5), \text{floor}(2.5) \rangle$$

$$= \langle -2, -1, 0, 1, 2 \rangle.$$

$$\text{map}(\text{floor} \circ \log_2, \langle 2, 3, 4, 5 \rangle)$$

$$= \langle \text{floor}(\log_2(2)), \text{floor}(\log_2(3)), \text{floor}(\log_2(4)), \text{floor}(\log_2(5)) \rangle$$

$$= \langle 1, 1, 2, 2 \rangle.$$

$$\text{map}(+, \langle (1, 2), (3, 4), (5, 6), (7, 8), (9, 10) \rangle)$$

$$= \langle +(1, 2), +(3, 4), +(5, 6), +(7, 8), +(9, 10) \rangle$$

$$= \langle 3, 7, 11, 15, 19 \rangle$$

- The map function is an example of a *higher-order* function, which is any function that either has a function as an argument or has a function as a value. This is an important property that most good programming languages possess.

A List of Squares

- Suppose we want to compute sequences of squares of natural numbers, such as 0, 1, 4, 9, 16. In other words, we want to compute $f: \mathbb{N} \rightarrow \text{lists}(\mathbb{N})$ defined by $f(n) = \langle 0, 1, 4, \dots, n^2 \rangle$. We have two different ways:

- First way: define $s(x) = x*x$ and then construct a definition for f in terms of map , s , and seq as follows.

$$\begin{aligned} f(n) &= \langle 0, 1, 4, \dots, n^2 \rangle \\ &= \langle s(0), s(1), s(2), \dots, s(n) \rangle \\ &= \text{map}(s, \langle 0, 1, 2, \dots, n \rangle) \\ &= \text{map}(s, \text{seq}(n)). \end{aligned}$$

- Second way: construct a definition for f without using the function s that we defined for the first way.

$$\begin{aligned} f(n) &= \langle 0, 1, 4, \dots, n^2 \rangle \\ &= \langle 0*0, 1*1, 2*2, \dots, n*n \rangle \\ &= \text{map}(*, \langle (0, 0), (1, 1), (2, 2), \dots, (n, n) \rangle) \\ &= \text{map}(*, \text{pairs}(\langle 0, 1, 2, \dots, n \rangle, \langle 0, 1, 2, \dots, n \rangle)) \\ &= \text{map}(*, \text{pairs}(\text{seq}(n), \text{seq}(n))). \end{aligned}$$