

BOOLEAN ALGEBRA

INEL4205 - SPRING 2008

OPERACIONES BÁSICAS

- AND - SALIDA ES "1" SI TODAS LAS ENTRADAS SON "1"
- OR - SALIDA ES "1" SI ALGUNA ENTRADA ES "1"
- NOT (COMPLEMENTO) - SALIDA ES "1" SI LA ENTRADA ES "0", "0" SI LA ENTRADA ES "1"

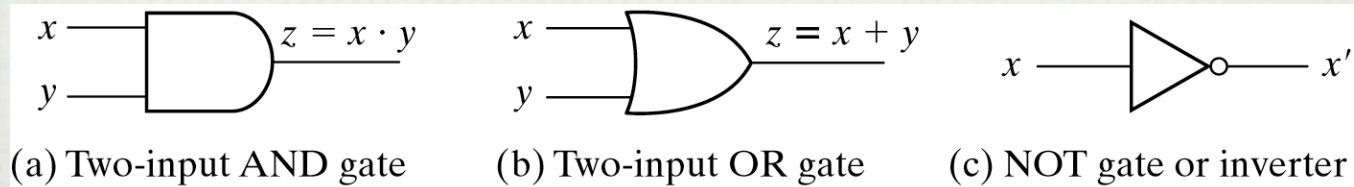


Fig. 1-4 Symbols for digital logic circuits

TABLAS DE VERDAD

- LA TABLA DE VERDAD PRESENTA LA SALIDA PARA TODAS LAS COMBINACIONES POSIBLES DE LAS ENTRADAS
- SI HAY "N" ENTRADAS, EL TOTAL DE COMBINACIONES UNICAS ES 2^N

A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

AND

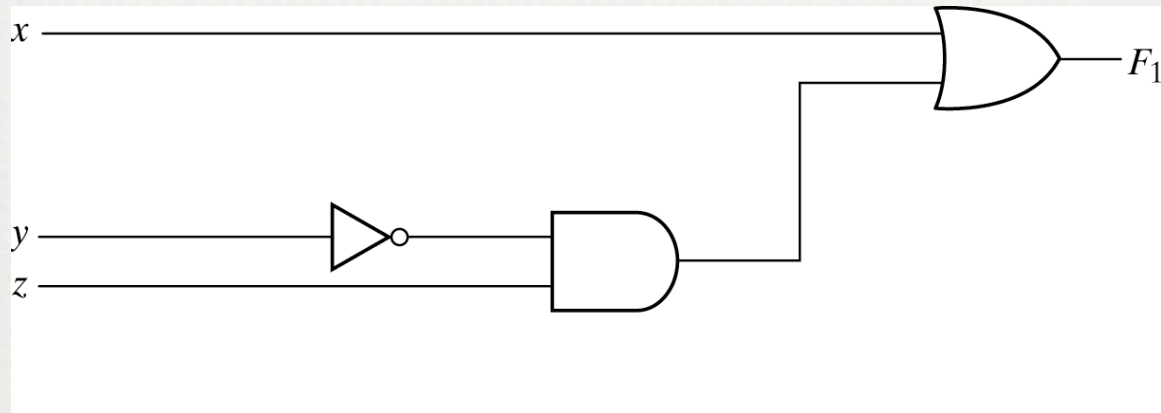
A	B	Y
0	0	0
0	1	1
1	0	1
1	1	1

OR

A	Y
0	1
1	0

NOT

EXPRESIÓN BOOLEANA



ESCRIBA LA TABLA DE VERDAD Y
LA EXPRESIÓN BOOLEANA QUE REPRESENTA
LA SALIDA " F_1 " EN TERMINOS DE LAS ENTRADAS
" x ", " y " Y " z "

Table 2-1*Postulates and Theorems of Boolean Algebra*

Postulate 2	(a)	$x + 0 = x$	(b)	$x \cdot 1 = x$
Postulate 5	(a)	$x + x' = 1$	(b)	$x \cdot x' = 0$
Theorem 1	(a)	$x + x = x$	(b)	$x \cdot x = x$
Theorem 2	(a)	$x + 1 = 1$	(b)	$x \cdot 0 = 0$
Theorem 3, involution		$(x')' = x$		
Postulate 3, commutative	(a)	$x + y = y + x$	(b)	$xy = yx$
Theorem 4, associative	(a)	$x + (y + z) = (x + y) + z$	(b)	$x(yz) = (xy)z$
Postulate 4, distributive	(a)	$x(y + z) = xy + xz$	(b)	$x + yz = (x + y)(x + z)$
Theorem 5, DeMorgan	(a)	$(x + y)' = x'y'$	(b)	$(xy)' = x' + y'$
Theorem 6, absorption	(a)	$x + xy = x$	(b)	$x(x + y) = x$

POSTULADO 4B FORMA 2

$$x + x'y = x + y$$

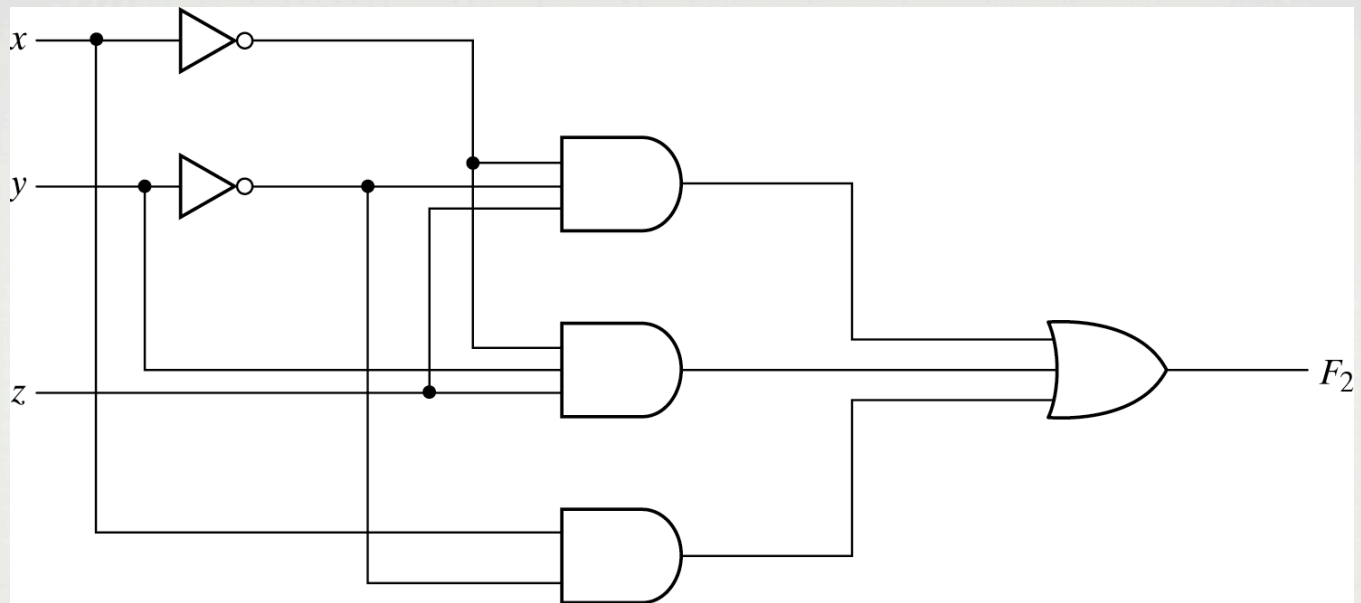
TEOREMA DEL CONSENSO

$$xy + x'z + yz = xy + x'z$$

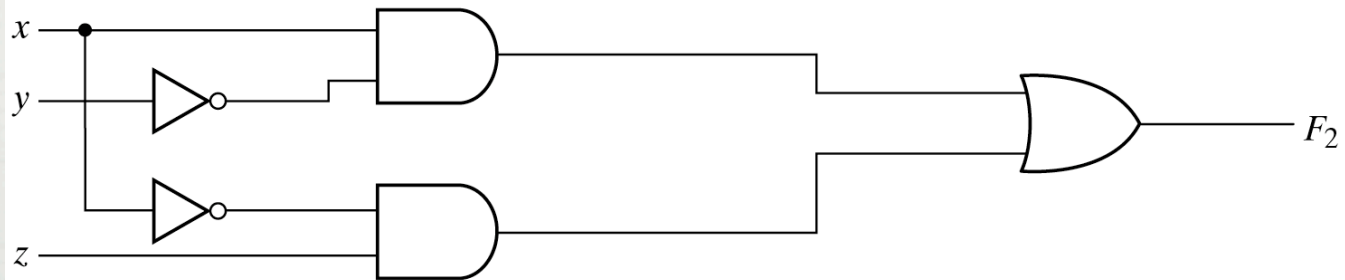
SIMPLIFICACIÓN

□ SIMPLIFIQUE LA SIGUIENTE EXPRESION BOOLEANA

$$F_2 = X'Y'Z + X'YZ + XY'$$



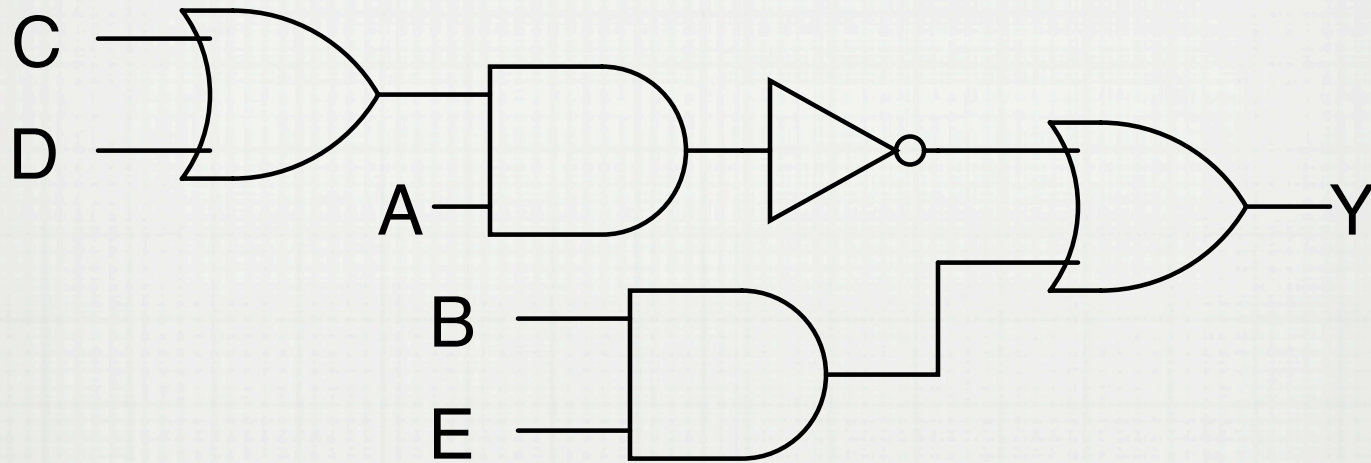
(a) $F_2 = x'y'z + x'yz + xy'$



(b) $F_2 = xy' + x'z$

Fig. 2-2 Implementation of Boolean function F_2 with gates

EJEMPLO



- ESCRIBA LA EXPRESIÓN BOOLEANA Y LA TABLA DE VERDAD PARA LA "Y" DEL DIAGRAMA DE ARRIBA
- DIBUJE EL DIAGRAMA PARA UN CIRCUITO QUE PRODUZCA LA SIGUIENTE SALIDA

$$ACF + DC'F'$$

EJEMPLOS

1. DETERMINE EL COMPLEMENTO DE LA SIGUIENTE EXPRESION BOOLEANA:

$$(A'+B)C'$$

2. DEMUESTRE QUE:

- $[(AB'+C)D'+E]' = [(A'+B)C'+D]E'$
- $X+XY=X$

EJEMPLOS

SIMPLIFIQUE:

- $x(x' + y)$
- $x + x'y$
- $(x + y)(x + y')$
- $xy + x'z + yz$
- $(x + y)(x' + z)(y + z)$

PRACTICA: PROBS. 2.2, 2.3 Y 2.4

TEOREMAS ADICIONALES

APLICANDO EL POSTULADO 4B

$$X + X'Y = X + Y$$

LLAMEMOSLO POSTULADO 4B FORMA 2

TEOREMA DEL CONSENSO

$$XY + X'Z + YZ = XY + X'Z + YZ(X + X')$$

$$XY + X'Z + YZ = XY(1 + Z) + X'Z(1 + Y)$$

$$XY + X'Z + YZ = XY + X'Z$$

SIMPLIFIQUE

$$ABCD' + A'B'CD + CD'$$

$$ABCD' + A'B'CD + CD'$$

$$(AB+1)CD' + A'B'CD$$

$$C(D' + A'B'D)$$

POR EL POSTULADO 4B FORMA 2

$$C(D' + A'B')$$

SIMPLIFIQUE

$$AB'C' + CD' + BC'D'$$

$$AB'C + CD' + BC'D'$$

$$AB'C + (C + BC')D'$$

POSTULADO 4B FORMA 2

$$AB'C + CD' + BD'$$

SIMPLIFIQUE

$$(A + B')(A' + B' + D)(B' + C + D')$$

$$(A + B')(A' + B' + D)(B' + C + D')$$

DUALIDAD

$$AB' + A'B'D + B'CD'$$

$$(A + A'D + CD')B'$$

$$(A + D + CD')B'$$

$$(A + D + C)B'$$

DUALIDAD

$$ACD + B'$$

$$XY + X'YZ' + YZ$$

$$XY + X'YZ' + YZ$$

$$XY + X'YZ' + YZ(x + x')$$

$$XY(1 + Z) + X'Y(Z' + Z)$$

$$XY + X'Y$$

$$Y$$

$$(xY' + Z)(x + Y') Z$$

$$(xY' + Z)(x + Y')Z$$

$$(xY' + Zx + ZY')Z$$

$$xY'Z + xZ + Y'Z$$

$$(x + 1)Y'Z + xZ$$

$$(x + Y')Z$$

$$xY' + Z + (x' + Y)Z'$$

$$XY' + Z + (X' + Y)Z'$$

POSTULADO 4B FORMA 2

$$XY' + X' + Y + Z$$

POSTULADO 4B FORMA 2

$$X' + Y' + Y + Z$$

1

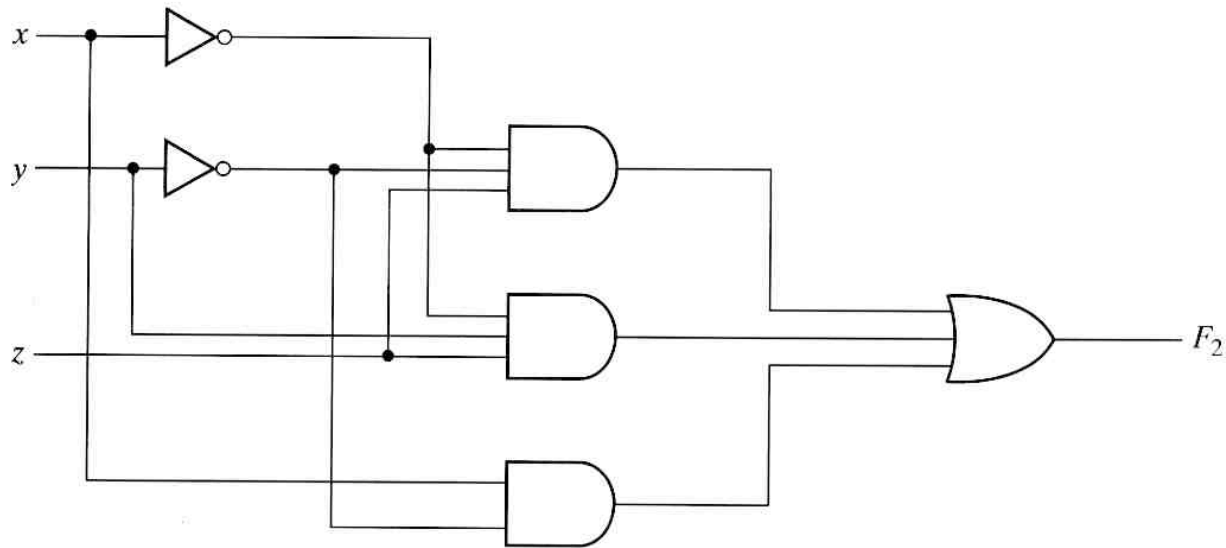
OTRA FORMA:

$$(XY' + Z)' = (X' + Y)Z'$$

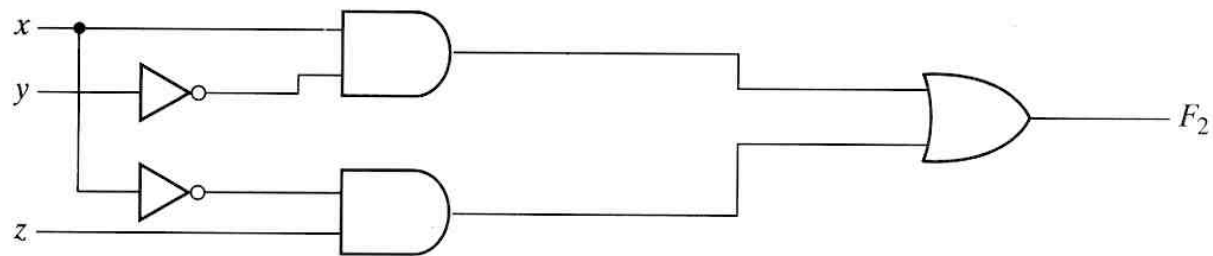
ASI QUE LA EXPRESIÓN TIENE LA FORMA

$$A + A'$$

Y POR LO TANTO SIEMPRE ES 1



(a) $F_2 = x'y'z + x'yz + xy'$ ← 3 TERMS; 8 LITERALS



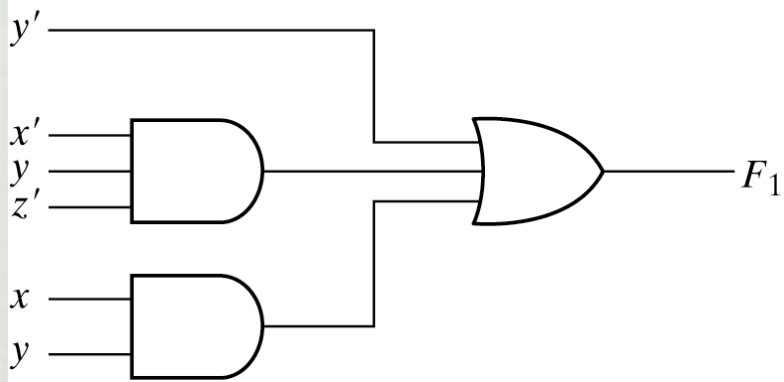
(b) $F_2 = xy' + x'z$ ← 2 TERMS; 4 LITERALS

FIGURE 2-2

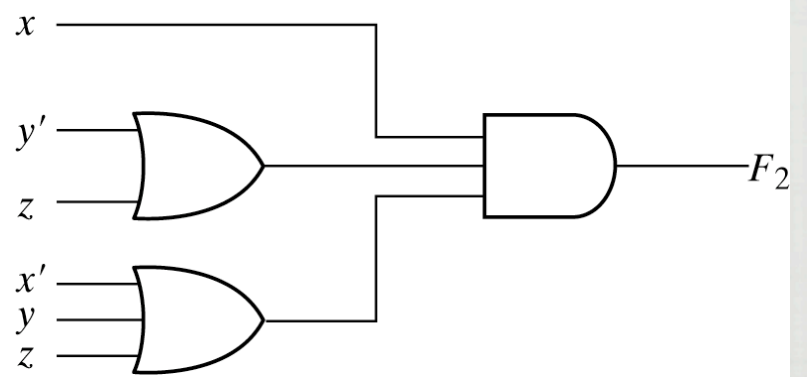
Implementation of Boolean function F_2 with gates

Table 2-3
Minterms and Maxterms for Three Binary Variables

x	y	z	Minterms		Maxterms	
			Term	Designation	Term	Designation
0	0	0	$x'y'z'$	m_0	$x + y + z$	M_0
0	0	1	$x'y'z$	m_1	$x + y + z'$	M_1
0	1	0	$x'yz'$	m_2	$x + y' + z$	M_2
0	1	1	$x'yz$	m_3	$x + y' + z'$	M_3
1	0	0	$xy'z'$	m_4	$x' + y + z$	M_4
1	0	1	$xy'z$	m_5	$x' + y + z'$	M_5
1	1	0	xyz'	m_6	$x' + y' + z$	M_6
1	1	1	xyz	m_7	$x' + y' + z'$	M_7

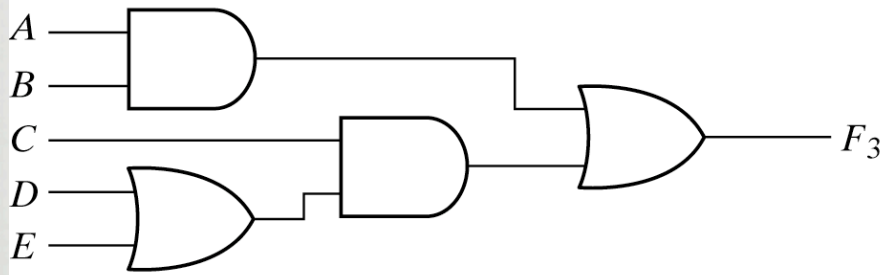


(a) Sum of Products

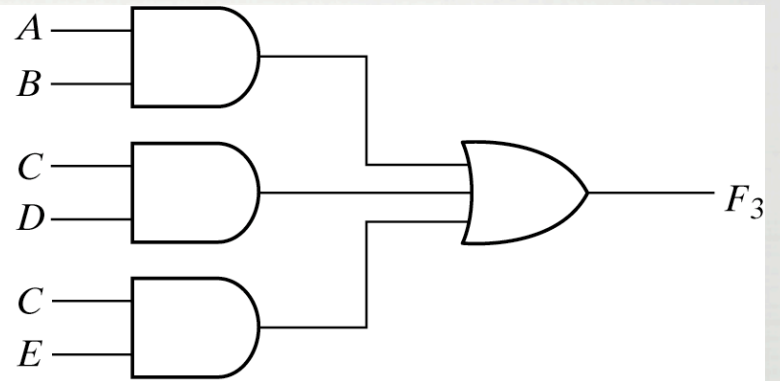


(b) Product of Sums

Fig. 2-3 Two-level implementation



(a) $AB + C(D + E)$



(b) $AB + CD + CE$

Fig. 2-4 Three- and Two-Level implementation

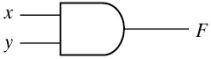
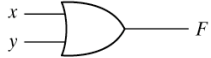
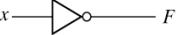
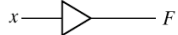

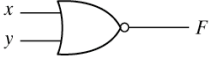


Name	Graphic symbol	Algebraic function	Truth table															
AND		$F = xy$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	0	1	0	0	1	1	1
x	y	F																
0	0	0																
0	1	0																
1	0	0																
1	1	1																
OR		$F = x + y$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	1	1	0	1	1	1	1
x	y	F																
0	0	0																
0	1	1																
1	0	1																
1	1	1																
Inverter		$F = x'$	<table border="1"> <thead> <tr> <th>x</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>1</td></tr> <tr><td>1</td><td>0</td></tr> </tbody> </table>	x	F	0	1	1	0									
x	F																	
0	1																	
1	0																	
Buffer		$F = x$	<table border="1"> <thead> <tr> <th>x</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td></tr> </tbody> </table>	x	F	0	0	1	1									
x	F																	
0	0																	
1	1																	
NAND		$F = (xy)'$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	1	1	0	1	1	1	0
x	y	F																
0	0	1																
0	1	1																
1	0	1																
1	1	0																
NOR		$F = (x + y)'$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	0	1	0	0	1	1	0
x	y	F																
0	0	1																
0	1	0																
1	0	0																
1	1	0																
Exclusive-OR (XOR)		$F = xy' + x'y$ $= x \oplus y$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>0</td></tr> <tr><td>0</td><td>1</td><td>1</td></tr> <tr><td>1</td><td>0</td><td>1</td></tr> <tr><td>1</td><td>1</td><td>0</td></tr> </tbody> </table>	x	y	F	0	0	0	0	1	1	1	0	1	1	1	0
x	y	F																
0	0	0																
0	1	1																
1	0	1																
1	1	0																
Exclusive-NOR or equivalence		$F = xy + x'y'$ $= (x \oplus y)'$	<table border="1"> <thead> <tr> <th>x</th> <th>y</th> <th>F</th> </tr> </thead> <tbody> <tr><td>0</td><td>0</td><td>1</td></tr> <tr><td>0</td><td>1</td><td>0</td></tr> <tr><td>1</td><td>0</td><td>0</td></tr> <tr><td>1</td><td>1</td><td>1</td></tr> </tbody> </table>	x	y	F	0	0	1	0	1	0	1	0	0	1	1	1
x	y	F																
0	0	1																
0	1	0																
1	0	0																
1	1	1																

Fig. 2-5 Digital logic gates

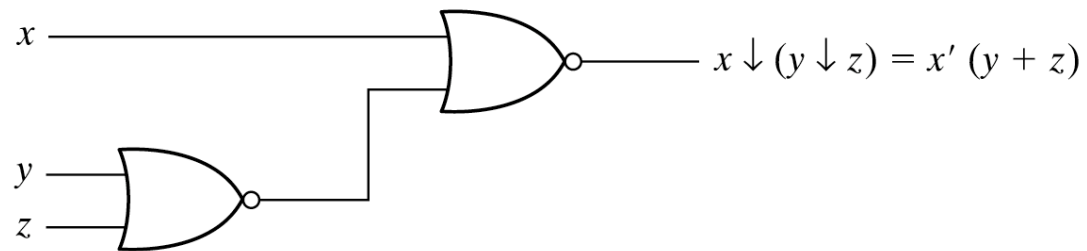
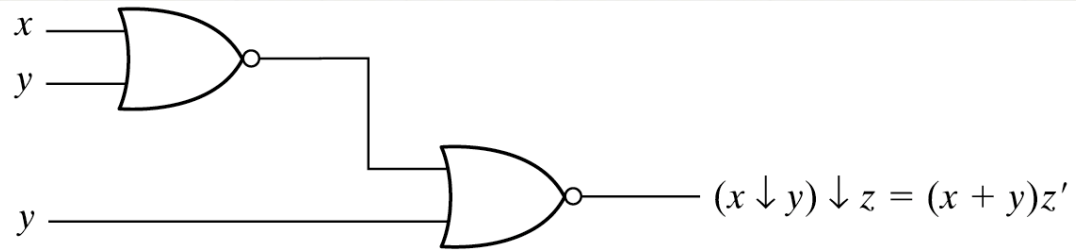
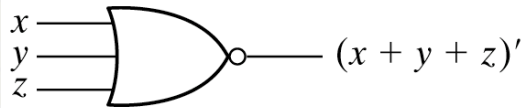
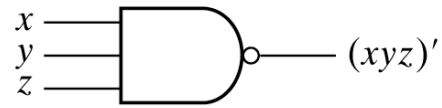


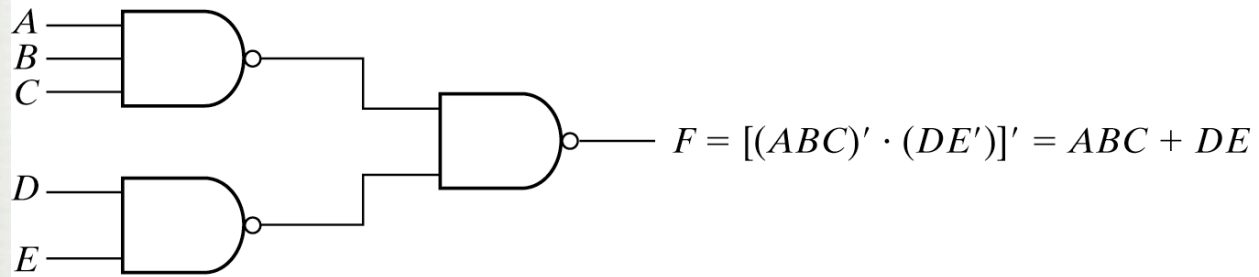
Fig. 2-6 Demonstrating the nonassociativity of the NOR operator; $(x \downarrow y) \downarrow z \neq x \downarrow (y \downarrow z)$



(a) 3-input NOR gate

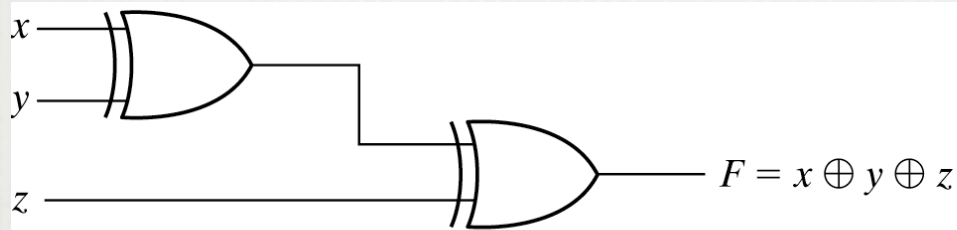


(b) 3-input NAND gate

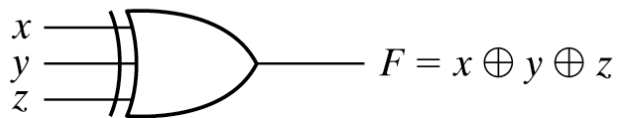


(c) Cascaded NAND gates

Fig. 2-7 Multiple-input and cascaded NOR and NAND gates



(a) Using 2-input gates



(b) 3-input gate

x	y	z	F
0	0	0	0
0	0	1	1
0	1	0	1
0	1	1	0
1	0	0	1
1	0	1	0
1	1	0	0
1	1	1	1

(c) Truth table

Fig. 2-8 3-input exclusive-OR gate

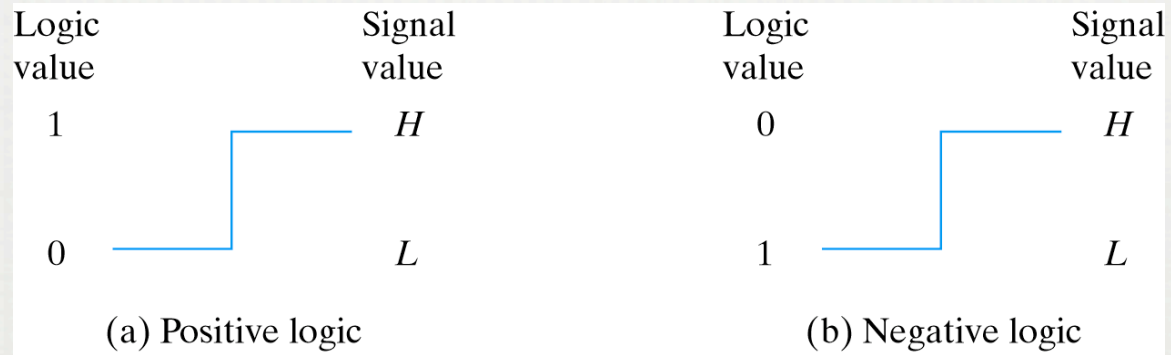
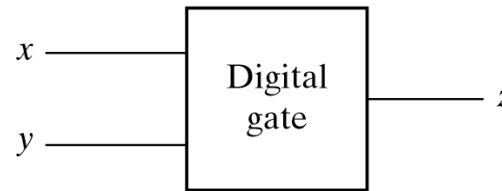


Fig. 2-9 signal assignment and logic polarity

x	y	F
L	L	L
L	H	L
H	L	L
H	H	H

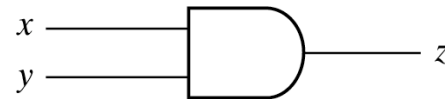
(a) Truth table with H and L



(b) Gate block diagram

x	y	z
0	0	0
0	1	0
1	0	0
1	1	1

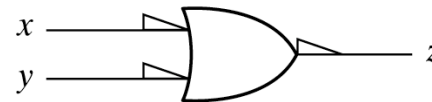
(c) Truth table for positive logic



(d) Positive logic AND gate

x	y	z
1	1	1
1	0	1
0	1	1
0	0	0

(e) Truth table for negative logic



(f) Negative logic OR gate

Fig. 2-10 Demonstration of positive and negative logic